The Big - M Method

The Big - M method is another method of removing artificial variables from the basis. In this method we assign coefficients to artificial variables, undesirable from the objective function. If objective function Z is to be minimized, then a very large positive price (called penalty) is assigned to each artificial variable. Similarly, if Z is to be maximized, then a very large negative price (also called penalty) is assigned to each of these variables. The penalty will designated by

-M for a maximization problem and +M for a minimization problem, where

M > 0.

The following are steps of the Algorithm for solving LPP by the Big - M method;

- (i) Express the LPP in the standard form by adding slack variables, surplus variables and artificial variables. Assign a zero coefficient to both slack and surplus variables and a very large positive coefficient +M (for min. case) and -M (for max. case) to artificial variable in the objective function.
- (ii) The initial basic feasible solution is obtained by assigning zero value to original variables.
- (iii) Calculate the value of $C_j Z_j$ in last row of simplex table and examine these values.
 - If all $C_j Z_j \ge 0$ then the current basic feasible solution is optimal.
 - If for a column k, $C_k Z_k$ is most negative and all entries in this column are negative, then the problem has unbounded optimal solution.
 - If one or more $C_j Z_j < 0$ (minimization case), then select the variable to enter into the basis with the largest negative $C_j Z_j$ value. That is $C_k Z_k = Min\{C_j Z_j\} : C_j Z_j < 0$.
- (iv) Determine the key row and key element in the same manner as discussed in the simplex algorithm of the maximization case.

Remarks

At any iteration of the simplex algorithm any one of the following cases may arise;

- 1. If at least one artificial variable is present in the basis with zero coefficient of *M* in each case $C_j Z_j \ge 0$, then the given LPP has no solution. That is, the current basic feasible solution is degenerate.
- 2. If at least one artificial variable is present in the basis with positive value and the coefficient of *M* in each $C_j Z_j \ge 0$, then given LPP has no optimum basic feasible solution. In this case the given LPP has a pseudo optimum basic feasible solution.

Example 1: Solve the following LPP using penalty (Big - M) method;

$$Max \ Z = x_1 + 2x_2 + 3x_3 - x_4$$

subject to

$$x_1 + 2x_2 + 3x_3 = 15$$

$$2x_1 + x_2 + 5x_3 = 20$$

$$x_1 + 2x_2 + x_3 + x_4 = 10$$

and

$$x_1, x_2, x_3 \ge 0$$

Solution

Since all constraints of the given LPP are equation, therefore adding only artificial variables A_1 and A_2 in the constraints. The standard form of the problem becomes;

$$Max \ Z = x_1 + 2x_2 + 3x_3 - x_4 - MA_1 - MA_2$$

subject to

$$x_1 + 2x_2 + 3x_3 + A_1 = 15$$

$$2x_1 + x_2 + 5x_3 + A_2 = 20$$

$$x_1 + 2x_2 + x_3 + x_4 = 10$$

and

 $x_1, x_2, x_3, A_1, A_2 \ge 0$

The initial basic feasible solution is given in Table1 below;

		$C_{\rm j} \rightarrow$	1	2	3	-1	-M	-M	
C_{B}	В	$b(=x_{\rm B})$	<i>x</i> ₁	<i>x</i> ₂	<i>x</i> ₃	<i>x</i> ₄	A_1	A_2	Min.Ratio
-M	A_1	15	1	2	3	0	1	0	$\frac{15}{3} = 5$
-M	A_2	20	2	1	Ģ	0	0	1	$\frac{20^{5}}{5} = 4 \rightarrow$
-1	<i>x</i> ₄	10	1	2	1	1	0	0	$\frac{10}{T} = 10$
Z = -35M - 10		$Z_{j} =$	-3M-1	-3M-2	-8M-1	-1	-M	-M	
		$C_{\rm j} - Z_{\rm j}$	3M+2	3M+4	8M+4	0	0	0	
					↑				

Table 1: Initial Solution

Since the value of $C_3 - Z_3$ in Table 1 has largest positive value the variable x_3 is chosen to enter into the basis. To get an improved basic feasible solution, apply the following row operations and removing A_2 from the basis.

 $\frac{R_2(new)}{\longrightarrow} \frac{R_2(old)}{5(key \ element)} = (4, \ 2/5, \ 1/5, \ 1, \ 0, \ 0)$ $R_1(new) \rightarrow R_1(old) - (3)R_2(new) = (3, -1/5, 7/5, 0, 0, 1)$ $R_3(new) \rightarrow R_3(old) - (1)R_1(new) = (6, \ 3/5, \ 9/5, 0, \ 1, \ 0)$

The improved solution is shown in Table 2

		$C_{\rm j} \longrightarrow$	1	2	3	-1	-M	
C_{B}	В	$b(=x_{\rm B})$	x_1	<i>x</i> ₂	<i>x</i> ₃	<i>x</i> ₄	A_1	Min.Ratio
-M	A_1	3	-1/5	7/5	0	0	1	$\frac{3}{7/5} = 15/7 \rightarrow$
3	<i>x</i> ₃	4	2/5	1/5	1	0	0	$\frac{4}{1/5} = 20$
-1	<i>x</i> ₄	6	3/5	9/5	0	1	0	$\frac{6}{9/5} = 30/9$
Z = -3M + 6		$Z_{j} =$	(M/5)-3/5	-(7M/5)-6/5	3	-1	-M	
		$C_{\rm j} - Z_{\rm j}$	-(M/5)-2/5	(7M/5)+16/5	0	0	0	
				↑				

Table 2: Improved Solution

The solution shown in Table 2 is not optimal because $C_2 - Z_2$ is positive. Thus, applying the following row operations for entering variable x_2 into the basis and removing variable A_1 from the basis.

$$R_1(new) \rightarrow \frac{R_1(old)}{7/5(keyelement)} = (15/7, -1/7, 1, 0, 0)$$

$$R_2(new) \rightarrow R_2(old) - (1/5)R_1(new) = (25/7, 3/7, 0, 1, 0)$$

$$R_3(new) \rightarrow R_3(old) - (9/5)R_1(new) = (15/7, 6/5, 0, 0, 1)$$

The new solution is shown in Table 3

Table 3: Improved Solution)n
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		$C_{\rm j} \rightarrow$	1	2	3	-1	
$C_{\rm B}$	В	$b(=x_{\rm B})$	x_1	x_2	<i>x</i> ₃	x_4	Min.Ratio
2	<i>x</i> ₂	15/7	-1/7	1	0	0	-
3	<i>x</i> ₃	25/7	3/7	0	1	0	$\frac{25/7}{3/7} = 25/3$
-1	<i>x</i> ₄	15/7	\bigcirc	0	0	1	$\frac{15/7}{6/7} = 5/2 \rightarrow$
Z = 90/7		$Z_{ m j}$	1/7	2	3	-1	
		$C_{\rm j} - Z_{\rm j}$	6/7	0	0	0	
			↑				

Again, the solution shown in Table 3 is not optimal. Thus, applying the following row operations by entering x_1 into the basis and removing variable x_4 from the basis.

 $\begin{array}{l} R_{3}(new) & \frac{R_{3}(old)}{6/7(key\ element)} = (15/6,\ 1,\ 0,\ 0,\ 7/6) \end{array}$

$$R_2(new) \to R_2(old) - (3/7)R_3(new) = (15/6, 0, 0, 1, -1/2)$$

$$R_1(new) \to R_1(old) - (-1/7)R_3(new) = (15/6, 0, 1, 0, 1/6)$$

The new solution is shown in Table 4

		$\stackrel{C_{j}}{\dashrightarrow}$	1	2	3	-1
C _B	В	$b(=x_{\rm B})$	<i>x</i> ₁	<i>x</i> ₂	<i>x</i> ₃	<i>x</i> ₄
2	x_2	15/6	0	1	0	1/6
3	<i>x</i> ₃	15/6	0	0	1	-1/2
1	x_1	15/6	1	0	0	7/6
Z =		Z_{j}	1	2	3	0
15						
		$C_{\rm j} - Z_{\rm j}$	0	0	0	-1

Table 4: Optimal Solution

Since all $C_j - Z_j \le 0$ in Table 4, Thus, an optimal solution has been arrived with values of variables as $x_1 = 15/6$, $x_2 = 15/6$, $x_3 = 15/6$, $x_4 = 0$ and *Max* Z = 15.

Example 2: Solve the following LPP using penalty (Big - M) method;

Main $Z = 600x_1 + 500x_2$

subject to

$$2x_1 + x_2 \ge 80 x_1 + 2x_2 \ge 60$$

and

 $x_1, x_2 \ge 0$

Solution:

By introducing surplus variables S_1 and S_2 and artificial variables A_1 and A_2 in the constraints. The standard form of the problem becomes;

Main $Z = 600x_1 + 500x_2 + 0S_1 + 0S_2 + MA_1 + MA_2$

subject to

$$2x_1 + x_2 - S_1 + A_1 = 80$$

$$x_1 + 2x_2 - S_2 + A_2 = 60$$

and

$$x_1$$
, x_2 , S_1 , S_2 , A_1 , $A_2 \ge 0$

The initial basic feasible solution is obtained by setting $x_1 = x_2 = S_1 = S_2 = 0$ as shown in Table 1;

Μ 600 500 Μ $C_i \rightarrow$ 0 0 В $b(=x_{\rm B})$ x_2 S_1 S_2 A_1 A_2 Min.Ratio $C_{\rm B}$ x_1 80 = 80М A_1 80 2 -1 0 1 0 60 $= 30 \rightarrow$ Μ A_2 60 1 0 0 -1 1 2 Z = 140M3M 3M -M Z_{i} -M Μ М 600-3M 500-3M Μ Μ 0 0 $C_i - Z_i$ ↑

Table 1: Initial Solution

Since the value of $C_2 - Z_2$ in Table 1 has largest negative value, therefore enter variable x_2 to replace basic variable A_2 into the basis. To get an improved basic feasible solution, apply the following row operations.

$$\frac{R_2(new)}{\rightarrow} \frac{\frac{R_2(old)}{2(key \, element)}}{2(key \, element)} = (30, 1/2, 1, 0, -1/2, 0)$$

$$R_1(new) \rightarrow R_1(old) - (1)R_2(new) = (50, 3/2, 0, -1, 1/2, 1)$$

The improved solution is shown in Table 2

		$C_{\rm j} \rightarrow$	600	500	0	0	М	
C_{B}	В	$b(=x_{\rm B})$	x_1	x_2	S_1	S_2	A_1	Min.Ratio
М	A_1	50	3/2	0	-1	1/2	1	$\frac{50}{3/2} = 33.33 \rightarrow$
500	<i>x</i> ₂	30	1/2	1	0	-1/2	0	$\frac{30}{1/2} = 60$
Z = 15000 + 50M		$Z_{\rm j}$	(3M/2)+250	500	-M	(M/2)-250	М	
		$C_j - Z_j$	350-3M	0	М	250-M/2	0	
			↑					

Table 2: Improved Solution

The solution shown in Table 2 is not optimal because $C_1 - Z_1$ is largest negative. Thus, applying the following row operations by entering variable x_1 into the basis and removing variable A_1 from the basis.

$$\begin{array}{ccc}
R_1(new) & R_1(old) \\
\xrightarrow{} & 3/2(key \ element) \\
R_2(new) & \rightarrow R_2(old) - (1/2)R_1(new) = (40/3, \ 0, \ 1, \ 1/3, \ -2/3) \\
\end{array}$$

The new solution is shown in Table 3

		$C_{\rm j} \rightarrow$	600	500	0	0
C_{B}	В	$b(=x_{\rm B})$	x_1	x_2	S_1	S_2
600	x_1	100/3	1	0	-2/3	1/3
500	<i>x</i> ₂	40/3	0	1	1/3	-2/3
Z = 80, 000/3		$Z_{ m j}$	600	500	-700/3	-400/3
		$C_{\rm j} - Z_{\rm j}$	0	0	700/3	400/3

Table 3 Optimal Solution

In Table 3 all the numbers in the $C_j - Z_j$ row are either zero or positive and also both artificial variables have been reduced to zero, an optimum solution has been arrived at with $x_1 = 100/3$, $x_2 = 40/3$ and total minimum cost, Z = 80, 000/3.